



**Logic in Practice**

- Horn Clauses
- Semantic Networks
- Description Logic
- Satisfiability**
- Logics of Action

# Logic in Practice

# Horn Clauses

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$$x \wedge y \rightarrow z \equiv \neg x \vee \neg y \vee z$$

Horn clause: at most one positive literal (exactly one is definite clause)

$\text{Cat}(x) \rightarrow \text{Furry}(x), \text{Meows}(x).$

$\text{Cat}(y) \rightarrow \text{Feline}(y).$

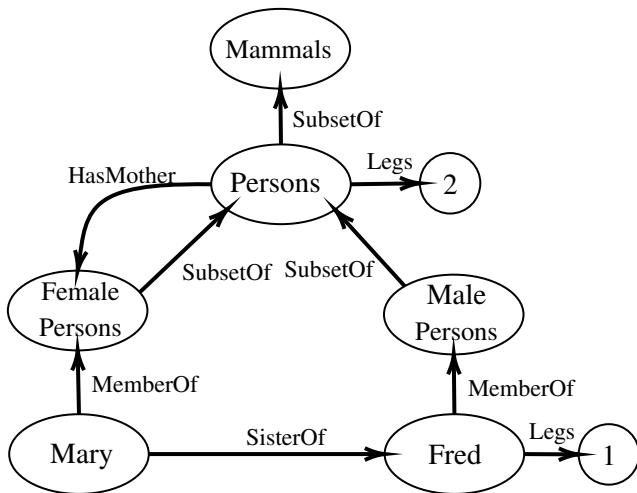
$\text{Furry}(A).$

$\text{Meows}(A).$

?  $\text{Cat}(z).$

Undecidable query

# Semantic Networks



# Semantic Networks

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Multiple Aspects:

- Visual notation
- Restricted Logic
- Set of implementation tricks

Typically in implementation:

- Efficient indexing
- Precomputation
- Methods for defaults

Also known as: frames, inheritance networks, semantic graphs, description logics, ontologies

# Description Logic

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Computing categories and membership, including:

- 1 subsumption
- 2 classification
- 3 inheritance

Missing many things:

- negation
- disjunction
- nested functions
- existentials
- intractability

# Description Logic

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- concepts (primitive and derived), instances
- roles and properties
- subsumptions:  $subsumes(x,y)$  iff:
  - $x$  is a concept and
  - $x, y$  have same ancestor and
  - for each role of  $x$  with restriction  $r_x$ ,  $y$  has the same role with a restriction  $r_y$  that  $r_x$  subsumes

# Satisfiability

# Boolean Satisfiability

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Given a boolean logic formula, is there a T/F assignment to all variables that make the formula true?

$$(a \vee b \vee c) \wedge (\neg a \vee b \vee \neg c) \wedge (\neg a \vee \neg b \vee c) \wedge (\neg a \vee \neg b \vee \neg c)$$



# Davis-Putnam-Logemann-Loveland Algorithm

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## DPLL( $\alpha$ )

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- 1: if  $\alpha$  has no clauses, **return** true
  - 2: if  $\alpha$  has an empty clause, **return** false
  - 3: if  $\alpha$  contains a unit clause, **return** DPLL(SIMPLIFY( $\alpha$ , *literal*))
  - 4:  $v \leftarrow$  choose a literal in  $\alpha$
  - 5: if DPLL(SIMPLIFY( $\alpha$ ,  $v$ )) is true, **return** true
  - 6: else **return** DPLL(SIMPLIFY( $\alpha$ ,  $\neg v$ ))
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## SIMPLIFY( $\alpha$ , *literal*)

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- 1: remove clauses in  $\alpha$  where *literal* is positive
  - 2: remove  $\neg$ *literal* from clauses where it appears
  - 3: **return** new  $\alpha$
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# Local SAT search

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- 1 Start with random solution
- 2 Repeatedly flip a variable to satisfy the most clauses
- 3 After a while, restart

GSAT, WalkSAT choose next variable differently

Both are much more efficient in practice than DPLL



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Practice

Satisfiability

**Logics of  
Action**

Event Calculus

Situation Calculus

Problems

# Logics of Action

# Event Calculus

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Events and fluents are reified:

$T(\text{Equals}(\text{President}(\text{USA}), \text{George}), \text{Begin}(1790), \text{End}(1790))$

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# Situation Calculus

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World state (situation) is reified:

$$\mathit{Result}(\mathit{GoForward}, s_0) = s_1$$

$$\mathit{Result}(\mathit{Turn}(\mathit{Left}), s_1) = s_2$$

$$\forall s, a, b \mathit{Clear}(a, s) \wedge \mathit{Clear}(b, s) \rightarrow \\ \mathit{On}(a, b, \mathit{Result}(\mathit{PutOn}(a, b), s))$$



# Problems with Logic

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- **Defaults:** difficult to build coherent semantics and efficient inference (default logics, probabilistic logic)
- **Ramification problem:** choosing what to infer (specialized systems)
- **Retraction:** when something changes from true to false (truth maintenance systems)
- **Qualification Problem:** making rules correct (probabilistic logic)

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Practice

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Action

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